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DIGITAL IMAGE REGISTRATION METHOD BASED UPON BINARY BOUNDARY MAPS

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TABLE OF CONTENTS

LIST OF ILLUSTRATIONS

LIST OF TABLES

FOREWORD

This study was initiated under the direction of Dr. R.R. Jayroe, Flight Data Statistics Office, Aerospace Environment Division, Aero-Astrodynamics Laboratory, Marshall Space Flight Center. This work was accomplished via the in-house efforts of Dr. R.R. Jayroe and Mr. C.W. Campbell of the same office, and by the efforts of Dr. J.F. Andrus of Northrop Services, Inc., Huntsville, Alabama, under contract NAS8-21810. Dr. Andrus must be credited for conceiving the original idea used in the binary correlation routine, and has continued to contribute to this study since joining the faculty of Louisiana State University, New Orleans, in August 1973.

DIGITAL IMAGE REGISTRATION METHODS BASED UPON BINARY BOUNDARY MAPS

I. INTRODUCTION

The ability to register or match the ground scene of images in the form of digital data serves an important role in the analysis of remotely sensed, multispectral, earth observation data and change detection. In the case of multispectral camera data, the images from the different camera stations must be registered before a proper analysis can be performed. Registration is also necessary when the data are acquired from several different sensors. In the case of change detection, where data are acquired from the same ground scene but at different time intervals, registration has an important role in determining what changes in shape have taken place in the ground scene, and registration permits an analysis to determine how the multispectral signatures of various ground features change as a function of time. For meteorological applications, registration methods based on meteorological satellite imagery can be used to estimate cloud velocities for global weather information.

Because of the typically large volumes of data involved, the use of binary maps in registration is recommended for the reasons discussed below. Converting the raw data to a binary boundary map typically represents a data compression of 60 to 70 percent of the original data. An additional significant compression of data is realized by working with sequences of, rather than individual, boundary points, which requires only the knowledge of the start scan and column and the length of the sequence.

If $A_1(x,y)$ and $A_2(x+\xi, y+\zeta)$ represent the amplitudes of the data in images 1 and 2 respectively, at scan coordinates x and $x+\xi$ respectively, and at column coordinates y and $y+\zeta$ respectively, then the average product $\overline{A_1(x,y)\,A_2(x+\xi, y+\zeta)}$ is computed over x and y for various combinations of ξ and ζ to determine the scan and column shifts necessary to register the two digital images. If this average is applied to the raw data, it may also be necessary to remove mean values and normalize in order to produce a correlation coefficient with a well defined peak. The use of the raw data appears to have some drawbacks as compared to using binary border maps. First, an ambiguity can result, since a large negative correlation peak can occur and represent a good match as well as a large positive peak. This is in most part due to the variation of the data acquired under different environmental conditions; i.e., different sun angles, seasons, atmospheric conditions, etc. On the other hand, binary boundary maps are produced from relative changes in the data, and the boundaries indicated in the ground scene tend not to change with environmental conditions. Since the binary boundary maps contain data that consist only of 0's and 1's, the average product $\overline{A_1(x,y)\,A_2(x+\xi, y+\zeta)}$ will always be positive and tends to be sharp so that removal of means and normalization is not necessary. In addition, it is possible to compute the average product $\overline{A_1(x,y)\,A_2(x+\xi, y+\zeta)}$ without multiplying when binary data are used. Addition can be used to replace multiplication, and this reduces computer computation time significantly.

Section II is a discussion of the production of binary boundary maps, while Sections III and IV are concerned with the correlation and registration schemes, respectively. Section V contains the results, test cases, and summary

II. BINARY BOUNDARY MAPS

The purpose of the binary boundary map is to categorize the digital data representing the ground scene into homogeneous areas and boundaries, and an example of a boundary map from ERTS digital imagery with some feature identification is shown in Figure 1.

The computer program used to generate the boundary map typically runs about 13 minutes (IBM-7094 time) on an area of 1000 scans by 255 columns, which is approximately 327 samples/second. This time includes the amount of time required to read and compile the program in addition to the calculation time. At present, no concerted effort has been made to optimize the running time of this program or to specialize its output for the registration program, but the efforts are in a holding status for future consideration.

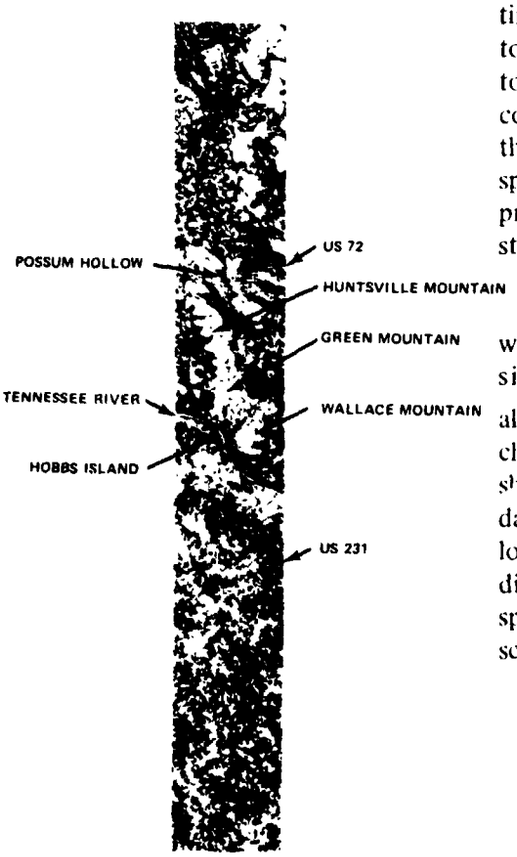

Figure 1. Boundary map from a portion of ERTS imagery.

The logic of the program is to work with two scans of raw digital data simultaneously. Let ${}_k x_{ij}$ represent the algebraic value of the data in the kth channel of data at scan i and column j as shown in Figure 2. If there are n channels of data, then k ranges from 1 to n. For each location i, j, the average vector component distances, which represent various degrees of spectroscopic changes within the ground scene, are computed using the formulas

$$S_x = \left[\frac{1}{n} \sum_{k=1}^{n} ({}_k x_{ij} - {}_k x_{i-1\,j})^2 \right]^{1/2} \qquad (1)$$

and

$$S_y = \left[\frac{1}{n} \sum_{k=1}^{n} (k^x_{i,j} - k^x_{i,j-1})^2 \right]^{1/2} , \qquad (1)$$

(Concluded)

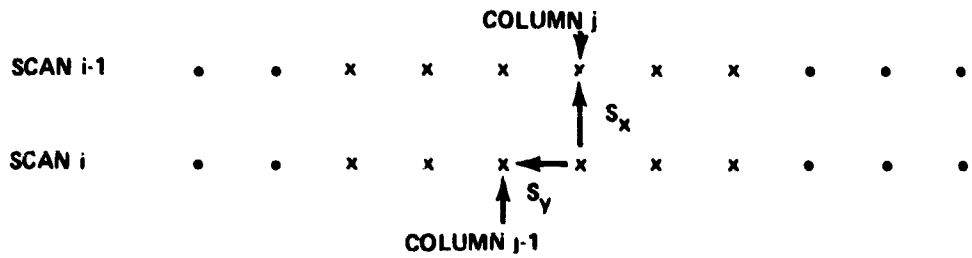

Figure 2. Computation of S_x and S_y for channel K.

and are stored in a joint histogram, $P(S_x, S_y)$. Dividing by n in equation (1) assures that S_x and S_y have the same typical range regardless of the number of data channels. Based upon previous experience with various types of data, it was found that the size of the joint histogram array could be limited in size to a 50 by 50 array such that any data element with a value of S_x or S_y larger than 50 could be automatically assigned as a boundary element. As the raw data are read into the program and the number of occurrences of S_x and S_y are accumulated in the joint histogram, a new tape is created for later use. This tape contains the numbers 1 through $50^2 = 2500$ and instead of writing S_x and S_y on this tape for each data element, one number is written on the tape that gives the location of S_x and S_y in the joint histogram. The location, I, of S_x and S_y in the joint histogram is computed using

$$I = (51)\, S_x + S_y \quad , \qquad (2)$$

and is a unique one-dimensional representation of the two dimensions S_x and S_y. The new tape containing a number I for each data element eliminates the necessity for recalculating S_x and S_y for each data element a second time. After all the data elements from the data set have been exhausted and the joint histogram is complete, a decision is made as to which combinations of S_x and S_y are to be considered as boundaries. The decision curve for a boundary element is based upon the formula

$$\left(\frac{S_x}{S_x' + ASCN}\right)^{IPOW} + \left(\frac{S_y}{S_y' + ACOL}\right)^{IPOW} > (2)BLIM \quad , \qquad (3)$$

where S_x' and S_y' are the values of S_x and S_y at the mode of the joint histogram and ASCN, ACOL, IPOW, and BLIM are input parameters to the program. The input parameters allow for a wide variety of decision curve shapes and positions. Nominally IPOW = 2 and BLIM = 1, whereas ASCN is usually equal to ACOL and must be estimated based upon experience. The possible values of I or correspondingly S_x and S_y are then inserted into equation (2) and the variable N(I) in the computer program is set equal to 0 if equation (2) is not satisfied and is set equal to 1 if equation (2) is satisfied. The tape containing I for each data element is then read into the program, and the value of N(I) is written out as another tape which is called the binary boundary map and contains only the numbers 0 and 1. The use of the intermediate tape containing the I values permits the boundary decision curve to be varied without recalculating S_x and S_y for every data element.

III. CORRELATION SCHEME

The development work for the correlation scheme was initiated because of a lack of a readily available registration program and after a review and discussion of some of the most recently published papers on registration, such as References 1 through 3. A computer listing of the correlation scheme is provided in Appendix A. The central idea is to compute only those correlation delays which are affected by overlapping boundary elements for each shift of the window against the picture.

Before describing the correlation scheme, it is necessary to define the computer program variables and terminology for the computer routine.

Picture	A section of a boundary map to be used as the reference data.
Window	A section of a boundary map which is to be registered with the picture. The window has one-half as many scans and columns of data as the picture.
NROW	The number of data scans in the window.
NCOL	The number of data samples per scan or columns of data in the window. Nominally, NROW is equal to NCOL.

$\frac{NCOR(I,J)}{(NROW)(NCOL)}$	The average product $\overline{A_1(x,y)\ A_2(x+I, y+J)}$. The value of NCOR(I,J) is the number of coinciding boundary points that occur when the upper left corner element of the window is placed upon the Ith row and Jth column of the picture. (I = 1,2, ..., NROW + 1 and J = 1, 2, ..., NCOL + 1.)
KMAX	Total number of boundary sequences in the window.
LMAXI	Total number of boundary sequences in the picture for I = 1, 2, ..., NROW.
LMAX	Total number of boundary sequences in the picture for all I, I = 1, 2, ..., (2) NROW.
NW(J)	Dummy variable for reading binary boundary map data. (J = 1, ..., NROW for window; J = 1, ..., (2)NROW for picture.)
K	Number of the Kth window boundary sequence (K = 1, ..., KMAX).
L	Number of the Lth picture boundary sequence (L = 1, ..., LMAX).
IW(K)	The value of IW(K) is the scan number minus one of the Kth boundary sequence in the window.
JW(K)	The value of JW(K) is the column number minus one of the beginning of the Kth boundary sequence in the window.
NSQW(K)	The value of NSQW(K) is the number of boundary elements minus one in the Kth boundary column sequence in the window.
IP(L)	The value of IP(L) is the scan number of the Lth boundary sequence in the picture.
JP(L)	The value of JP(L) is the column number of the beginning of the Lth boundary sequence in the picture.
NSQP(L)	The value of NSQP(L) is the number of boundary elements minus one in the Lth boundary column sequence in the picture.
I	The value of I is the scan shift in NCOR(I,J).

JL	The value of JL is the lower value of the column shift J in NCOR(I,J).
JU	The value of JU is the upper value of the column shift J in NCOR(I,J). (J = JL, JL + 1, ..., JU.)
NAD	The value of NAD is the number of coinciding boundary points at a given I and J for the Kth window boundary sequence and the Lth picture boundary sequence.
JM	The value of JM controls the value of NAD: JM causes NAD to decrease as the Kth and Lth boundary sequences decreasingly overlap because of an increase in J from JL to JU.
NU	For a given I, the value of NU is the column number of the end of the Lth boundary sequence in the picture.
NSPAN	For a given I, the value of NSPAN is the column number of the end of the Kth boundary sequence in the window.
MINL	MINL is a dummy variable which changes and represents the smaller of the window or picture boundary sequence minus one.
CALL TIMNOW ( )	Subroutine for calling internal clock to time computation of NCOR(I,J).

The first step of the program is to read in the sections of the binary boundary maps to be considered, skipping all of the nonboundary or zero elements, and storing the start scan, start column, and column length of the sequences in the window and picture. This information is stored consistent with the definition of IW(K), JW(K), NSQW(K), IP(L), JP(L), NSQP(L), KMAX, LMAX1, and LMAX, and results in a considerable compression of the data. The calculation of the above variables is accomplished in the DO 20 nested do-loop for the window and the DO 40 nested do-loop for the picture. The remaining part of the program contains two major calculation segments, the DO 100 nested do-loop and the DO 200 nested do-loop. Because of their similarity it will suffice to explain only one of these do-loops and then point out their differences. The DO 100 segment is concerned with the entire window and the top half of the picture; i.e., I = 1, 2, ..., NROW. The two outer do-loops, DO 100 and DO 80, determine the I and J shifts that cause the individual boundaries of the K window sequences to coincide with the individual boundaries of the L picture sequences. No negative shifts are permitted; i.e., $1 \leq I \leq NCOL + 1$ and $1 \leq J \leq NCOL + 1$. Within the DO 100 and DO 80 loops, there are 3 cases to be considered. The first case is the DO 56 loop and it occurs when the minimum length, MINL, of either the window or picture sequence is one boundary element. This occurrence is illustrated in Figure 3 and can result in an addition of at most 1 to NCOR(I,J) for each value of J. The case which probably occurs most often is

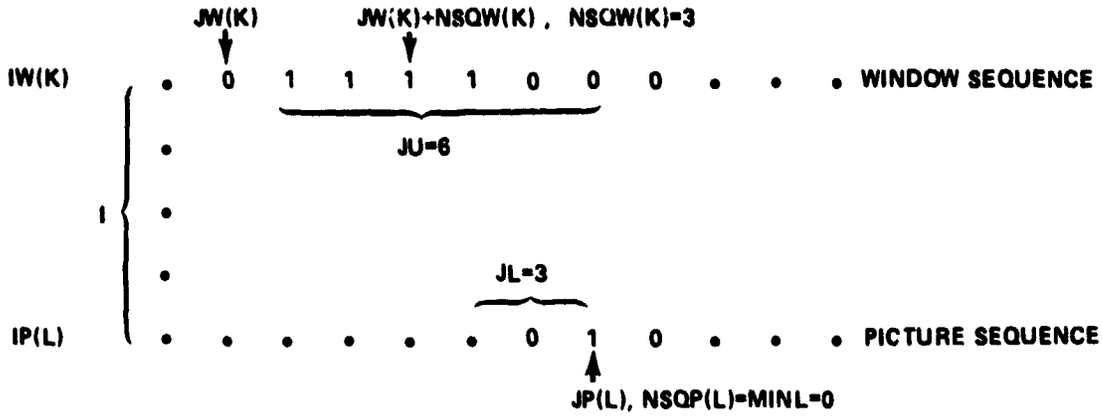

Figure 3. Subcase for DO 56 loop.

the DO 68 loop, and a typical example is shown in Figure 4. The first time through the DO 68 loop NAD = 1 is added to NCOR(I,J) and NAD continues to increase by one until it is greater than MINL, the minimum length of either sequence minus one. Once NAD is one greater than MINL it remains that value until J is greater than JM, and then NAD decreases by one each time until the two sequences no longer overlap in the column direction. Thus, both the initial and final values of NAD added to NCOR(I,J) are one, while the maximum value of NAD is the number of boundary elements in the smaller of the two sequences. Table 1 is a convenient way of viewing the number of occurrences of each column shift due to the presence of both sequences. The top two rows list the possible values of J or the column shift, while the next three rows list the possible values

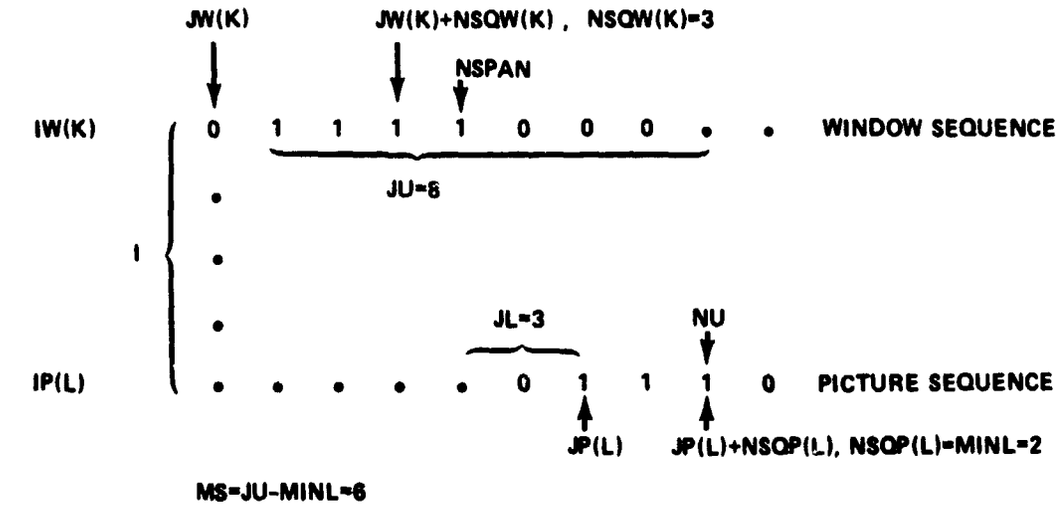

Figure 4. Subcase for DO 68 loop, JL > 0.

TABLE 1. NUMBER OF OCCURRENCES OF EACH COLUMN SHIFT FOR FIGURE 4

	JL		JM			JU
	3	4	5	6	7	8
JP(L)	JL	JL+1	JL+2	JL+3		
JP(L) + 1		JL+1	JL+2	JL+3	JL+4	
JP(L) + NSQP(L)			JL+2	JL+3	JL+4	JL+5
NAD	1	2	3	3	2	1

of J for each boundary element in the picture. The last row is the corresponding value for NAD for each J or the total number of occurrences of each column shift. An additional subcase for the DO 68 loop can occur when JL is negative, JP(L) > JW(K), and the length of the window sequence is equal to or less than the picture sequence. An example of this subcase is shown in Figure 5 which corresponds to Table 2.

The final two subcases occur in connection with the DO 78 loop. Both subcases have the minimum column shift, JL, which is initially negative. The first subcase has JP(L) less than or equal to JW(K) and is illustrated in Figure 6 which corresponds to Table 3. The second subcase is illustrated in Figure 7 which corresponds to Table 4.

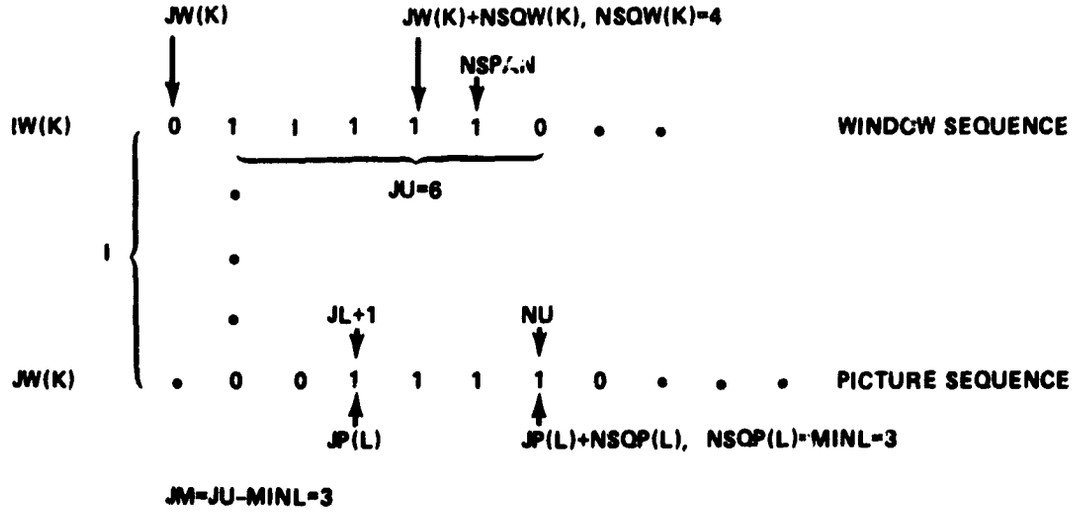

Figure 5. Subcase for DO 68 loop, JL ≤ 0, NSPAN ≤ NU, JP(L) > JW(K).

TABLE 2. NUMBER OF OCCURRENCES OF EACH COLUMN SHIFT FOR FIGURE 5

	JL	JM				JU
	1	2	3	4	5	6
JP(L)	JL	JL+1	JL+2			
JP(L) + 1	JL	JL+1	JL+2	JL+3		
JP(L) + 2	JL	JL+1	JL+2	JL+3	JL+4	
JP(L) + NSQP(L)		JL+1	JL+2	JL+3	JL+4	JL+5
NAD	3	4	4	3	2	1

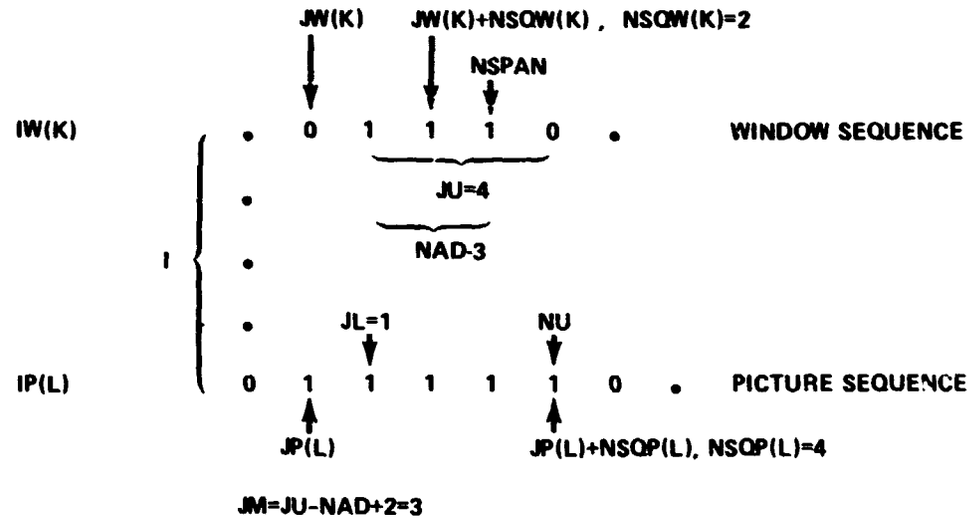

Figure 6. Subcase for DO 78 loop, JL ⩽ 0, JP(L) ⩽ JW(K).

TABLE 3. NUMBER OF OCCURRENCES OF EACH COLUMN SHIFT FOR FIGURE 6

	JL		JM	JU
	1	2	3	4
JP(L) + 1	JL			
JP(L) + 2	JL	JL+1		
JP(L) + 3	JL	JL+1	JL+2	
JP(L) + NSQP(L)		JL+1	JL+2	JL+3
NAD	3	3	2	1

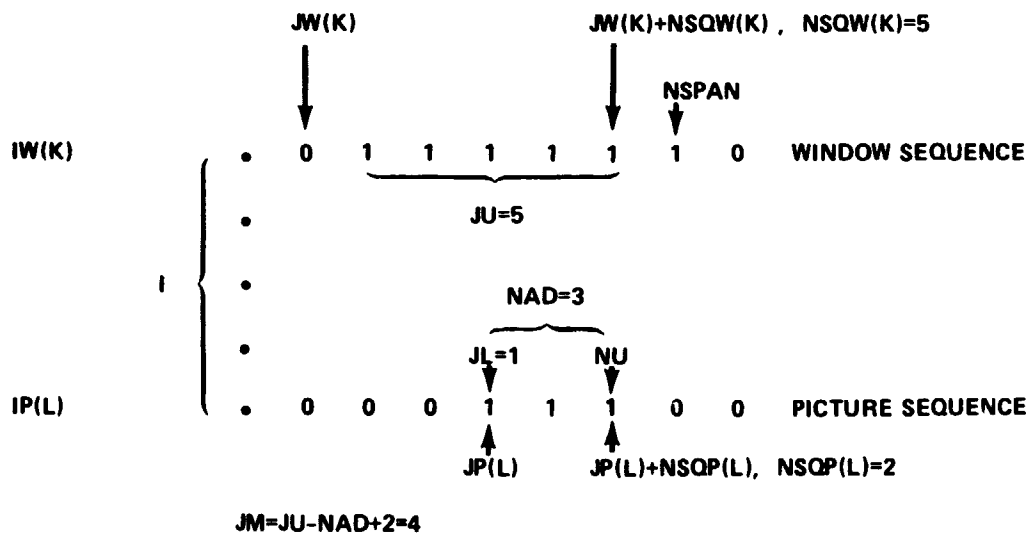

Figure 7. Subcase for DO 78 loop, JL $\leqslant$ 0, JP(L) > JW(K), NSPAN $\geqslant$ NU.

TABLE 4. NUMBER OF OCCURRENCES OF EACH COLUMN SHIFT FOR FIGURE 7

	JL			JM	JU
	1	2	3	4	5
JP(L)	JL	JL+1	JL+2		
JP(L) + 1	JL	JL+1	JL+2	JL+3	
JP(L) + NSQP(L)	JL	JL+1	JL+2	JL+3	JL+4
NAD	3	3	3	2	1

The second segment of the program or the DO 200 loop is identical to the first segment of the program, the DO 100 loop, except for the statement numbers. The main difference is that the DO 200 loop, while considering the entire window, only considers the bottom half of the picture. Thus, the DO 100 loop considers less and less of the first half of the picture sequences for all of the window sequences, and the DO 200 loop considers less and less of the window for the last half of all of the picture sequences. In the DO 100 loop, the scan shift I starts at the maximum available value and decreases to I = 1, while in the DO 200 loop I starts at the minimum available value and increases to a maximum of I = NROW + 1. A scan shift of I = 1 and a column shift of J = 1 indicate that the window and picture are already registered.

IV. REGISTRATION SCHEME

In applying the registration scheme it is assumed that the data can be misaligned by translation and/or rotation, but that no distortion exists between the two digital images. If significant distortions do exist over the entire images, then it may become necessary to register subportions of the digital images as separate data sets.

It is also assumed that both digital images are of the same scale. If the images are not of the same scale, then it is possible to scale down one image by locating several corresponding boundary elements on both images, choose a corresponding reference boundary element on each image, and compute for each image the average distance of all the other boundary elements from the reference boundary element. The ratio of the two average distances gives a scaling factor. The raw data image with the larger average distance can then be scaled to the same size as the other image, and a new boundary map can be computed.

To register two boundary maps, choose one map as the reference map or picture and choose subportions of the other map to be used as windows as illustrated in Figure 8. Figure 9 shows a portion of the two maps overlayed and the minimum and maximum scan and column shifts for the window. For the registration to work properly, the final

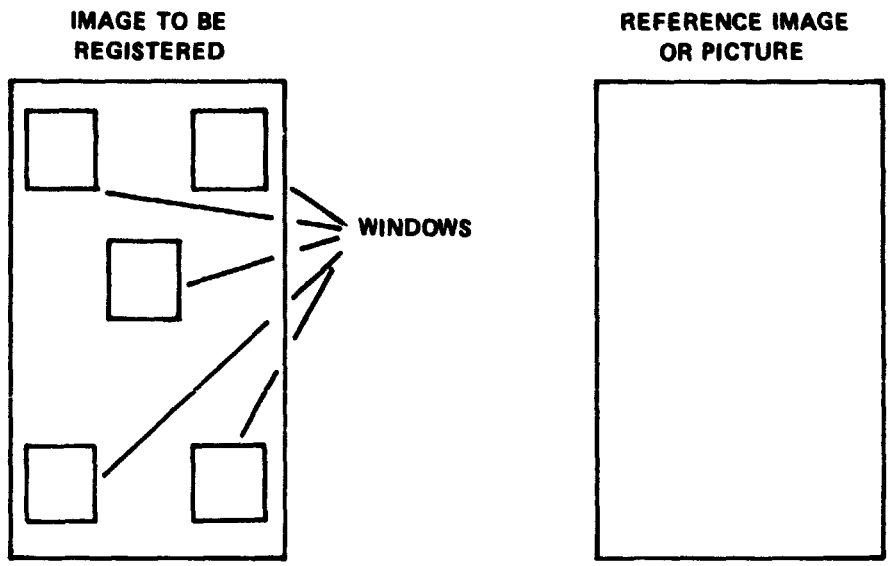

Figure 8. Window selection.

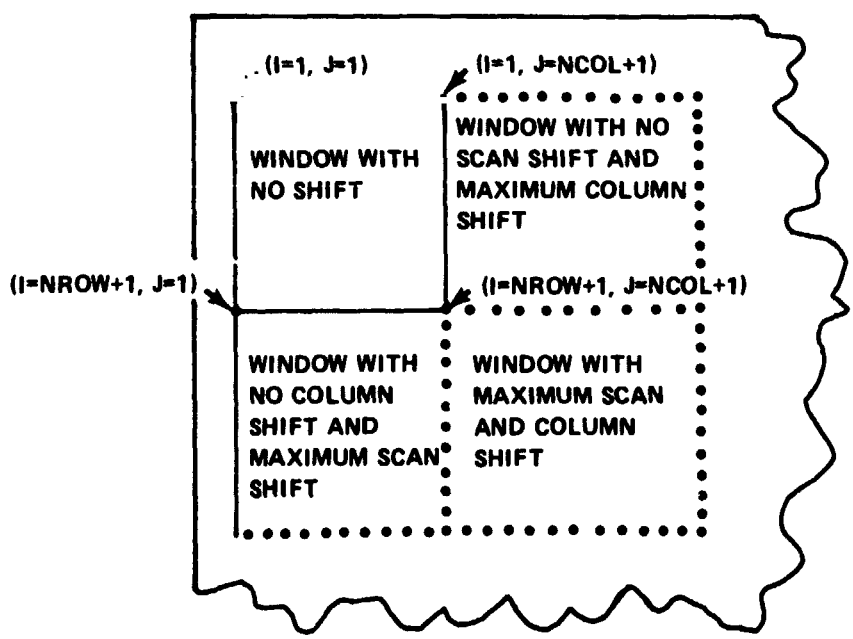

Figure 9. Overlay of digital images.

two assumptions are that the portions of the picture corresponding to the window are within the allowable range of the window and that the amount of rotation present in the picture or the window can be neglected as far as the computation of NCOR(I,J) is concerned and can be treated as a local translation. If the latter assumption is not valid, it may be necessary to repeat the registration procedure again after the data have been translated and/or rotated based upon the first registration effort. In the event that the first assumption is correct, it should be possible to overcome the second assumption by repeated application of the registration procedure to the data.

The mathematics of the translation-rotation required to register the two images is given below. Let ${}_wx_i$, ${}_wy_i$ be the scan and column coordinates respectively of the upper left corner on the ith window in the image to be registered, and let I_i and J_i correspond to the maximum value of NCOR(I,J) for the ith window. For any one of the windows, for example the jth, let the coordinates

$$ {}_px_j = {}_wx_j + I_j - 1 $$

and (4)

$$ {}_py_j = {}_wy_j + J_j - 1 $$

on the picture be designated as the center of rotation. Thus, if the entire window image is shifted by the amounts I_1-1 and J_1-1, the coordinates $({}_px_j, {}_py_j)$ and $({}_wx_j, {}_wy_j)$ will coincide and the rest of the image can be rotated about these coordinates to complete the registration, provided a rotation exists. If no rotation exists, the images are registered and the maximum value of NCOR(I,J) for the rest of the window occurs at I = J = 1. Assuming that rotation does exist, it will be necessary to compute the following quantities,

$$ \overline{{}_px{}_wx} = \frac{1}{N-1} \sum_{i \neq j}^{N} ({}_px_i - {}_px_j)({}_wx_i - {}_wx_j) \quad , \tag{5} $$

$$ \overline{{}_py{}_wy} = \frac{1}{N-1} \sum_{i \neq j}^{N} ({}_py_i - {}_py_j)({}_wy_i - {}_wy_j) \quad , \tag{6} $$

$$\overline{{}_{p}x_{w}y} = \frac{1}{N-1} \sum_{i \neq j}^{N} ({}_{p}x_i - {}_{p}x_j)({}_{w}y_i - {}_{w}y_j) \quad , \tag{7}$$

$$\overline{{}_{w}x_{p}y} = \frac{1}{N-1} \sum_{i \neq j}^{N} ({}_{w}x_i - {}_{w}x_j)({}_{p}y_i - {}_{p}y_j) \quad , \tag{8}$$

and

$$\tan\theta = -\frac{\overline{{}_{p}x_{w}y} - \overline{{}_{w}x_{p}y}}{\overline{{}_{p}x_{w}x} + \overline{{}_{p}y_{w}y}} \quad , \tag{9}$$

where N is the total number of windows. To fetch the data from the image to be registered, it is first necessary to solve for θ, and for each picture coordinate $({}_{p}x, {}_{p}y)$, it is necessary to solve the equations,

$${}_{p}x \cos\theta + {}_{p}y \sin\theta - {}_{p}x_j \cos\theta - {}_{p}y_j \sin\theta + {}_{p}x_j = {}_{w}x$$

and (10)

$$-{}_{p}x \sin\theta + {}_{p}y \cos\theta - {}_{p}x_j \sin\theta - {}_{p}y_j \cos\theta + {}_{p}y_j = {}_{w}y$$

for the appropriate unregistered image coordinates. In general, the coordinates $({}_{w}x, {}_{w}y)$ will not be integers, but in most cases it is probably most appropriate and less time consuming to round off to the nearest integer coordinate rather than try to interpolate the value of the data between the original data points.

The registration routine is currently being programmed and debugged by MSFC's Computation Laboratory, Engineering Computation Division (S&E-COMP-RRV). The routine has not yet been modified to accept the correlation scheme, but instead accepts a manual input of corresponding boundary elements from two different images.

V. TEST CASES, RESULTS, AND SUMMARY

The binary correlation routine was compared to a correlation and Fast Fourier Transform (FFT) correlation routine by obtaining running times from an internal clock on the IBM-7044 computer. A listing of the correlation routine and a program for using the FFT correlation routine are presented in Appendices B and C. The Fast Fourier Transform used, FFT7, is not presented in this report since it has been evaluated and listed in Reference 4 along with nine other FFT routines. According to Reference 4, FFT7 was the most versatile and ranked third in speed to the fastest routine, FFT3, which was twice as fast but worked only on real data, performed only a (+i) transform, and produced only a half transform. The second fastest transform was FFT8, which was only 25 milliseconds faster than FFT7.

The data set on which the programs were run is shown in Figure 10. In the data set, only the locations of the boundaries are printed out and represented as 1's, and the window and the picture both start in the upper left corner. The size of the data array used for the window is the first NCOL columns by NROW scans, and the picture is twice as large. For the Fast Fourier Transform the window has to be the same size as the picture, and consequently zeros had to be added to the window array. The window and picture data arrays were arranged such that they were already in registration or such that the maximum value of NCOR(I,J) occurred at I = J = 1.

The output of the binary correlation routine for a 32 by 32 window array is shown in Figure 11, and the maximum number of NCOR(I,J) is the total number of boundary elements in the window. Figure 12 is a graph of correlation lag array size versus running time for the binary correlation, Fast Fourier Transform correlation, and correlation routines. The actual running times for each program versus correlation array size are listed in Table 5. Storage requirements for the FFT correlation routine were such that it was not possible to compute 41 by 41 correlation lag points or larger, and the irregular running times as shown in Figure 12 are a result of the FFT routine being more efficient for some array sizes than others. The correlation routine was able to handle 57 by 57 correlation lag points, but the running time was already 5 minutes for 41 by 41 correlation lag points. Since it appears that running times for the correlation routine can easily be extrapolated from Figure 12 and Table 5, no larger size arrays were run.

According to Reference 1, the Sequential Similarity Detection Algorithm (SSDA) mentioned therein is approximately 50 times faster than the FFT correlation method. The factor of 50 is estimated from time ratios relating arithmetic operations on the IBM 360/65 used in computing the numerator of the correlation coefficient. By relating all arithmetic operations to integer adds for a 32 by 32 window and a 256 by 256 picture, the equivalent integer adds for the correlation, FFT correlation, and SSDA respectively were 2.57×10^8 and 2.2×10^6, which indicate a factor of 100 times faster than the FFT correlation. To compare the binary correlation with the SSDA it was necessary to put a counter in the DO loops containing NCOR(I,J) and to use the formula of

Figure 10. Binary boundary map data set.

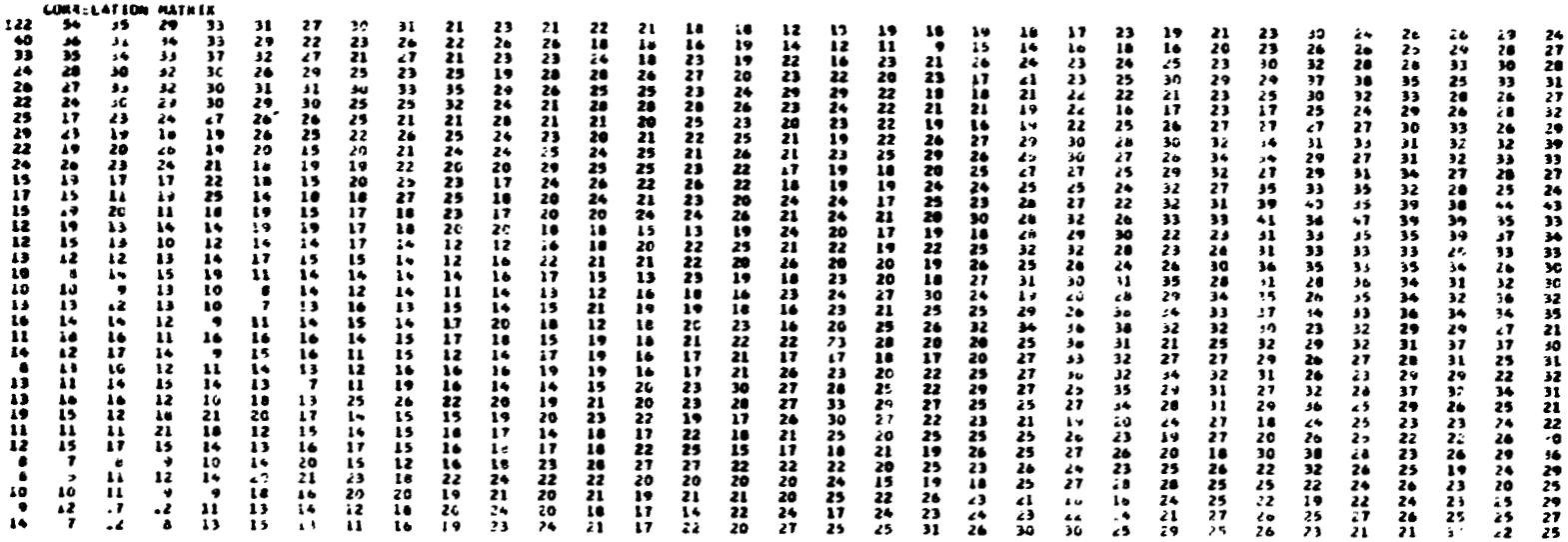

Figure 11. Binary correlation output for a 32 by 32 window array.

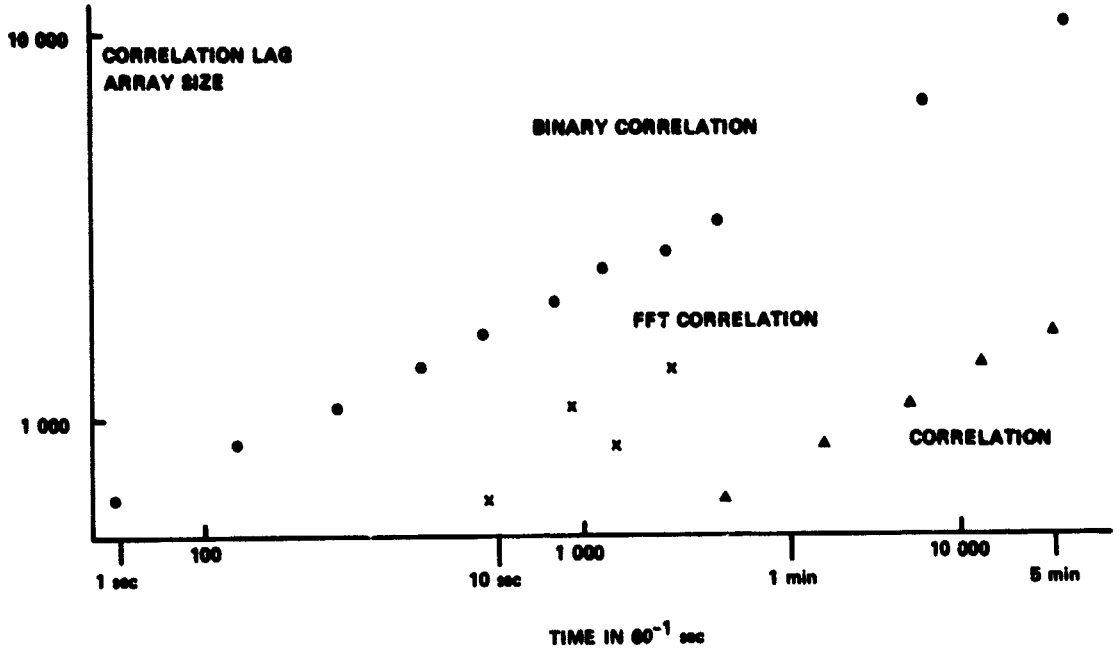

Figure 12. Correlation lag array size versus running time.

TABLE 5. RUNNING TIME VERSUS CORRELATION LAG ARRAY SIZE

Lag Array Size	Running Time in 60^{-1} seconds/Times Slower Than Binary Correlation		
	Binary Correlation	FFT Correlation	Correlation
25 by 25	58	585/10.09	2 450/42.24
29 by 29	124	1267/10.22	4 465/36.01
33 by 33	230	983/4.27	7 532/32.75
37 by 37	381	1778/4.67	11 962/31.4
41 by 41	559		18 097/32.37
45 by 45	850		
49 by 49	1 167		
53 by 53	1 710		
57 by 57	2 377		
81 by 81	8 371		
101 by 101	19 761		

Reference 1 for computing the equivalent number of integer adds for a 32 by 32 window and 64 by 64 picture. The formula presented in Reference 1 is $4(1 + 10\sqrt{M/32})(L - M+1)^2$ equivalent integer adds, where L and M are the picture and window length or width, respectively. For the above mentioned picture and window the number of equivalent integer adds for the SSDA is 47,916 while the number of integer adds for the particular data set shown in Figure 10 is 20,804. Thus, it would appear that the binary correlation could possibly be 2.3 times faster than the SSDA. However, it is also possible that computing equivalent integer adds for various types of programs does not give the complete story, since it is necessary in all the computer programs to execute additional computer statements other than updating the correlation. Otherwise, the binary correlation would be approximately 200 times faster than the FFT correlation. Table 5 indicates that it is not. It is realized that one does not obtain "something for nothing," and the processing time needed to produce a binary boundary map could negate the speed of the binary correlation. However, two points in addition to those mentioned in Section I are worth considering. According to Reference 1 it may become necessary to resort to the use of boundary maps when the channels under consideration are widely separated in spectral wavelength. This is certainly evidenced when one is working with near infrared or thermal imagery which tends to be quite noisy. Secondly, Figure 13 suggests a way to optimize the production of a boundary map.

Figure 13 illustrates the number of variables needed to define various window sizes for the correlation and binary correlation methods using the data set illustrated in Figure 10. Also included in Figure 13 is the number of variables needed to define the location of the boundary elements only in the data set, which is twice the number of boundary points. For the correlation it is necessary to store the entire window array, while for the binary correlation it is necessary only to store 3 X KMAX variables, the start scan, the start column, and the sequence length for each boundary sequence. Thus, Figure 13 indicates the amount of data compression that is possible. Table 6 lists the compression factor for various window sizes, which is the total number of data points divided by 3 X KMAX.

The production of a boundary map could be optimized by recognizing that three variables are needed to define each boundary sequence. Also, from the logic used in calculating IW(K), JW(K), NSQW(K), IP(L), JP(L), and NSQP(L), it is only necessary to determine boundaries within a scan and not necessary to determine boundaries that occur across two adjacent scan lines. Thus, the boundary program described in Section II could be optimized by working with one scan of data at a time instead of two and by eliminating all boundary sequences that contain less than three boundary elements. This would also optimize the compression by assuring that 3 X KMAX is always equal to or less than the number of boundary elements.

It is foreseen that in some cases it may become necessary to subtract mean values from NCOR(I,J) and normalize to produce a proper registration of two data sets. In this respect it is interesting to examine the correlation coefficient in terms of rewards and

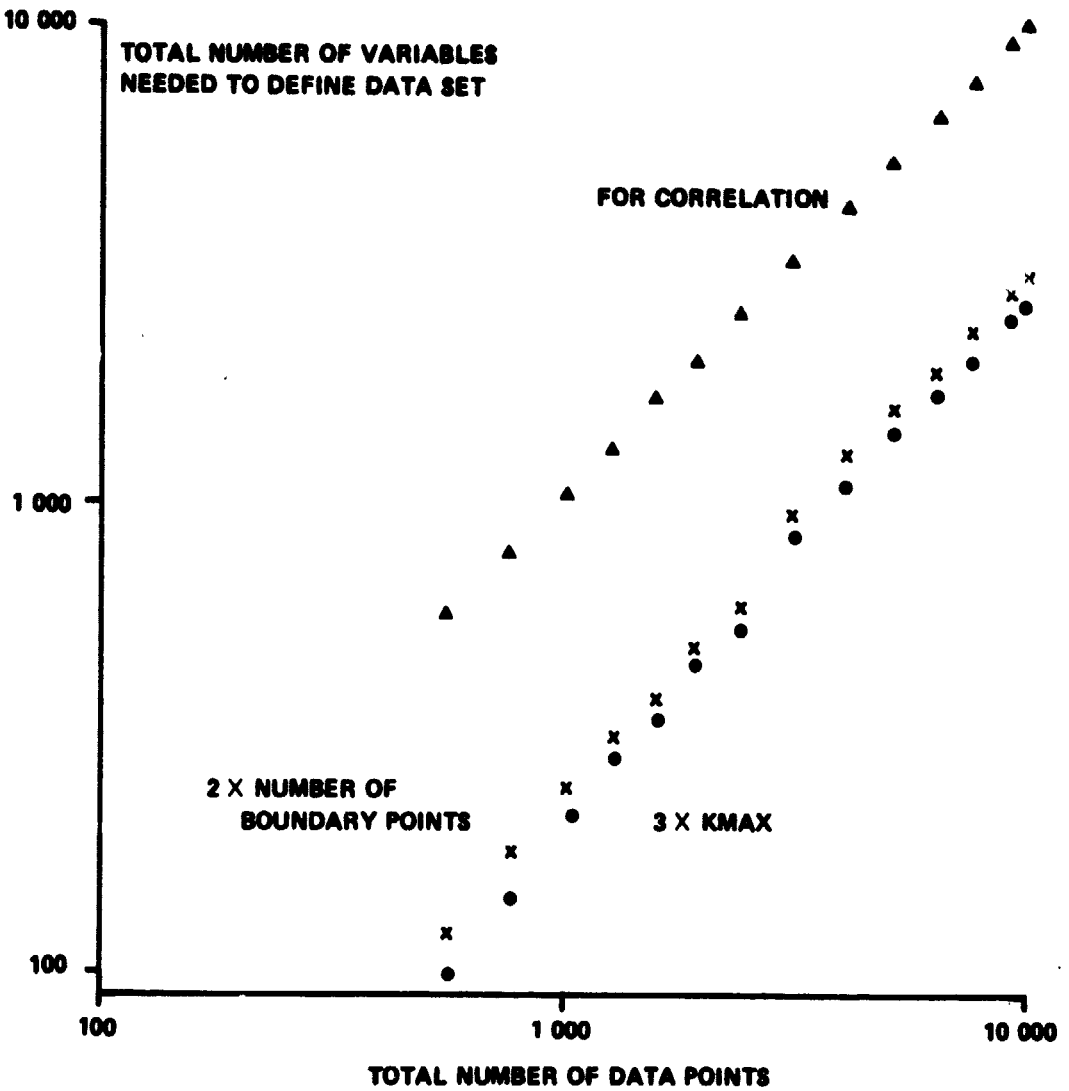

Figure 13. Number of variables needed to define window array.

penalties used in the binary correlation method. The binary correlation method presented in Section III only rewards for boundary elements that match and does not penalize for mismatches of boundary and nonboundary elements. By subtracting mean values and normalizing, it is possible to obtain rewards, for matching boundary elements and matching nonboundary elements, and penalties for mismatches of boundary and nonboundary elements. Before examining the correlation coefficient, it will be necessary to define a few terms.

TABLE 6. COMPRESSION FACTOR FOR WINDOW ARRAY

Window Array Size	Number of Boundary Elements	KMAX	Compression Factor
24 by 24	61	33	5.81
28 by 28	90	48	5.44
32 by 32	122	70	4.88
36 by 36	158	95	4.55
40 by 40	183	115	4.64
44 by 44	237	151	4.27
48 by 48	296	178	4.31
52 by 52	368	229	3.94
56 by 56	450	281	3.72
64 by 64		355	3.85
72 by 72		451	3.83
80 by 80	903	552	3.86
88 by 88		656	3.93
96 by 96		789	3.89
100 by 100	1431	847	3.94

NWB — The number of boundary elements in the window.

NPB(I,J) — The number of boundary elements in the picture that corresponds to the area covered by the window. NPB(I,J) changes as the window is moved to different I and J locations.

W — The mean value of the data in the window, W=NWB/[(NROW)(NCOL)].

P(I,J) — The mean value of the data in the picture that corresponds to the area covered by the window, P(I,J)=NPB(I,J)/[(NROW)(NCOL)].

The numerator of the correlation coefficient, C(I,J) can then be written as

$$C(I,J) = K_1 \cdot NCOR(I,J) + K_2 \cdot [NPB(I,J) - NCOR(I,J)]$$

$$+ K_3 \cdot [NWB - NCOR(I,J)] + K_4 \cdot [NROW \cdot NCOL - NPB(I,J)$$

$$+ NCOR(I,J) - NWB] \quad , \qquad (.1)$$

where

$$K_1 = [1 - W][1 - P(I,J)] \quad , \quad K_2 = -W[1 - P(I,J)] , \quad K_3 = -[1 - W]P(I,J)$$

and

$$K_4 = W \cdot P(I,J)$$

Since $0 \leq W, P(I,J) \leq 1$, then $K_1, K_4 \geq 0$ and $K_2, K_3 \leq 0$, and K_1 and K_4 are always greater than K_2 and K_3 in absolute magnitude. If the constants, K, are thought of in terms of rewards and penalties, then the first term of equation (11) is a reward times the number of boundary elements that match in the window and picture. The second and third terms, since they are negative, penalize the mismatch of boundary and nonboundary elements in the picture and window, respectively. The last term is a reward for nonboundary elements that match in the picture and window. It is interesting to note that for $W+P(I,J) < 1$, then $K_1 > K_4 > K_2, K_3$; that for $W+P(I,J)=1$, then $K_1 = K_4 > K_2, K_3$; and that for $W+P(I,J) > 1$, then $K_4 > K_1 > K_2, K_3$. This indicates that the binary correlation procedure gives a larger reward for matching up whatever occurs the least in the window or picture, boundaries or nonboundaries. This is also what a human observer would tend to do. To determine a normalization factor for equation (11), consider a perfect match for the window and picture; i.e., NWB=NPB(I,J)=NCOR(I,J) and $W=P(I,J) < 1$. Equation (11) becomes

$$C(I,J) = (1 - W)^2 \cdot NWB + W^2 \cdot (NROW \cdot NCOL - NWB)$$

$$= [1 - P(I,J)]^2 \cdot NPB(I,J) + [P(I,J)]^2$$

$$[NROW \cdot NCOL - NPB(I,J)] \qquad (12)$$

Thus, for a nonperfect match, the normalization factor should be some combination of the two expressions in equation (12). Since the two expressions in equation (12) are the variances of the window and picture respectively, the correct normalization factor would be the square root of the products of the two expressions to produce a correlation coefficient.

It does not appear that computing the correlation coefficient would add a significant amount of running time to the binary correlation routine. This is because the mean value and variance of the window need only be calculated once and $NWB=KMAX+\sum_{K=1}^{KMAX} NSQW(K)$. The calculation that would require additional time and storage is determining NPB(I,J) for the mean value and variance of the picture for each shift of the window. In most cases, however, it is anticipated that it will be necessary to compute only NCOR(I,J).

George C. Marshall Space Flight Center
National Aeronautics and Space Administration
Marshall Space Flight Center, Alabama, November 19, 1973

APPENDIX A. BINARY CORRELATION LISTING

```
5245, CAMPBELL      ,140145,30,                    FORTRAN SOURCE LIST
      ISN           SOURCE STATEMENT

        0 $IBFTC MAIN     DECK
        1       DIMENSION NW(200),FMT(12),NCOR(101,101),Iw(1324),Jw(1324)
        2       DIMENSION NSQW(1324),IP(3300),JP(3300),NSQP(3300)
        3       READ 300,NROW,NCOL
        6   300 FORMAT(2I5)
        7       PRINT 302,NROW,NCOL
       10   302 FORMAT(1X,5HNROW=,I4,1X,5HNCOL=,I4,/)
       11       READ(5,301)(FMT(L),L=1,12)
       16   301 FORMAT(12A6)
       17       NROW1=NROW+1
       20       NCOL1=NCOL+1
       21       NCOR(NROW1,NCOL1) =0
       22       K=0
       23       DO 20 I=1,NROW
       24       IN=-1
       25       READ(5,FMT)(NW(I1),I1=1,NCOL)
       32       DO 10 J=1,NCOL
       33       NCOR(I,J)=0
       34       IF(NW(J).EQ.0)GO TO 5
       37       IF(IN.GT.0) GO TO 3
       42       K=K+1
       43       IW(K)=I-1
       44       JW(K)=J-1
       45       NSQW(K)=0
       46       IN=1
       47       GO TO 10
       50     3 NSQW(K)=NSQW(K)+1
       51       GO TO 10
       52     5 IN=-1
       53    10 CONTINUE
       55    20 CONTINUE
       57       KMAX=K
       60       READ(5,301)(FMT(L),L=1,12)
       65       NROW2=NROW+NROW
       66       NCOL2=NCOL+NCOL
       67       L=0
       70       DO 40 I=1,NROW2
       71       IN=-1
       72       READ(5,FMT)(NW(I1),I1=1,NCOL2)
       77       DO 30 J=1,NCOL2
      100       IF(NW(J).EQ.0)GO TO 25
      103       IF(IN.GT.0) GO TO 23
      106       L=L+1
      107       NSQP(L)=0
      110       IP(L)=I
      111       JP(L)=J
      112       IN=1
      113       IF(I.GT.NROW1)GO TO 30
      116       LMAX1=L
      117       GO TO 30
      120    23 NSQP(L)=NSQP(L)+1
      121       GO TO 30
      122    25 IN=-1
      123    30 CONTINUE
      125    40 CONTINUE
```

```
5245, CAMPBELL      ,140145,30,                    FORTRAN SOURCE LIST MAIN
        ISN         SOURCE STATEMENT

        127         LMAX=L
        130         KSTRT=1
        131         ISUM=0
        132         CALL TIMNOW(IT1)
        133         DO 100 L=1,LMAX1
        134         NU=JP(L)+NSQP(L)
        135         DO 80 K=1,KMAX
        136         IF(IW(K).GE.IP(L))GO TO 100
        141         I=IP(L)-IW(K)
        142         IF(I.NE.1)GO TO 45
        145         IF(NU.GT.JW(K))GO TO 46
        150         GO TO 100
        151      45 IF(NU.LE.JW(K))GO TO 80
        154      46 JL=JP(L)-JW(K)-NSQW(K)
        155         IF(JL.GT.NCOL1)GO TO 80
        160         JU=NU-JW(K)
        161         IF(JL.LT.1)GO TO 70
        164         IF(NSQP(L).LE.NSQW(K))GO TO 53
        167         MINL=NSQW(K)
        170         GO TO 54
        171      53 MINL=NSQP(L)
        172      54 IF(MINL.GT.0)GO TO 60
        175         IF(JU.LE.NCOL1)GO TO 55
        200         JU=NCOL1
        201      55 DO 56 J=JL,JU
        202         ISUM=ISUM+1
        203      56 NCOR(I,J)=NCOR(I,J)+1
        205         GO TO 80
        206      60 JM=JU-MINL
        207         NAD=0
        210         IF(JU.LE.NCOL1)GO TO 64
        213         JU=NCOL1
        214      64 DO 68 J=JL,JU
        215         ISUM=ISUM+1
        216         IF(J.GT.JM)GO TO 66
        221         IF(NAD.GT.MINL)GO TO 68
        224         NAD=NAD+1
        225         GO TO 68
        226      66 NAD=NAD-1
        227      68 NCOR(I,J)=NCOR(I,J)+NAD
        231         GO TO 80
        232      70 NSPAN=JW(K)+NSQW(K)+1
        233         IF(JP(L).GT.JW(K))GO TO 71
        236         IF(NSPAN.GE.NU)GO TO 69
        241         NAD=NSQW(K)+1
        242         GO TO 75
        243      69 NAD=JU
        244         GO TO 75
        245      71 IF(NSPAN.GE.NU)GO TO 73
        250         NAD=NSPAN-JP(L)
        251         IF(NSQP(L).GT.NSQW(K))GO TO 74
        254         MINL=NSQP(L)
        255         GO TO 76
        256      74 MINL=NSQW(K)
        257      76 JM=JU-MINL
```

```
5245, CAMPBELL      ,140145,00,                    FORTRAN SOURCE LIST MAIN
      ISN       SOURCE STATEMENT

      260       JL=1
      261       GO TO 64
      262    75 NAD=NSQP(L)+1
      263    75 JM=JU-NAD+2
      264       DO 78 J=1,JU
      265       ISUM=ISUM+1
      266       IF(J.LT.JM)GO TO 78
      271       NAD=NAD-1
      272    78 NCOR(I,J)=NCOR(I,J)+NAD
      274    80 CONTINUE
      276   100 CONTINUE
      300       LMAX1=LMAX1+1
      301       DO 200 L=LMAX1,LMAX
      302       NU=JP(L)+NSQP(L)
      303    52 DO 180 K=KSTRT,KMAX
      304       I=IP(L)-IW(K)
      305       IF(I.GT.NROW1)GO TO 85
      310       IF(NU.LE.JW(K))GO TO 180
      313       JL=JP(L)-JW(K)-NSQW(K)
      314       IF(JL.GT.NCOL1)GO TO 180
      317       JU=NU-JW(K)
      320       IF(JL.LT.1)GO TO 701
      323       IF(NSQP(L).LE.NSQW(K))GO TO 531
      326       MINL=NSQW(K)
      327       GO TO 541
      330   531 MINL=NSQP(L)
      331   541 IF(MINL.GT.0)GO TO 601
      334       IF(JU.LE.NCOL1)GO TO 552
      337       JU=NCOL1
      340   552 DO 561 J=JL,JU
      341       ISUM=ISUM+1
      342   561 NCOR(I,J)=NCOR(I,J)+1
      344       GO TO 180
      345   601 JM=JU-MINL
      346       NAD=0
      347       IF(JU.LE.NCOL1)GO TO 641
      352       JU=NCOL1
      353   641 DO 681 J=JL,JU
      354       ISUM=ISUM+1
      355       IF(J.GT.JM)GO TO 661
      360       IF(NAD.GT.MINL)GO TO 681
      363       NAD=NAD+1
      364       GO TO 681
      365   661 NAD=NAD-1
      366   681 NCOR(I,J)=NCOR(I,J)+NAD
      370       GO TO 180
      371   701 NSPAN=JW(K)+NSQW(K)+1
      372       IF(JP(L).GT.JW(K))GO TO 711
      375       IF(NSPAN.GE.NU)GO TO 691
      400       NAD=NSQW(K)+1
      401       GO TO 751
      402   691 NAD=JU
      403       GO TO 751
      404   711 IF(NSPAN.GE.NU)GO TO 731
      407       NAD=NSPAN-JP(L)
```

```
5245, CAMPBELL     ,140145,00,                FORTRAN SOURCE LIST MAIN
      ISN          SOURCE STATEMENT

       410         IF(NSQP(L).GT.NSQW(K))GO TO 741
       413         MINL=NSQP(L)
       414         GO TO 761
       415   741   MINL=NSQW(K)
       416   761   JM=JU-MINL
       417         JL=1
       420         GO TO 641
       421   731   NAD=NSQP(L)+1
       422   751   JM=JU-NAD+2
       423         DO 781 J=1,JU
       424         ISUM=ISUM+1
       425         IF(J.LT.JM)GO TO 781
       430         NAD=NAD-1
       431   781   NCOR(I,J)=NCOR(I,J)+NAD
       433   180   CONTINUE
       435         GO TO 200
       436    85   KSTRT=K+1
       437         GO TO 52
       440   200   CONTINUE
       442         CALL TIMNOW(IT2)
       443         PRINT 5555,ISUM
       444  5555   FORMAT(5X,7HISUM = ,I7)
       445         JTIME=IT2-IT1
       446         PRINT 318
       447   318   FORMAT(5X,50H**************************************************)
       450         PRINT 317,JTIME
       451   317   FORMAT(/////,5X,15HELAPSED TIME = ,I5,11H/60 SECONDS,//////)
       452         PRINT 1100,KMAX,LMAX
       453  1100   FORMAT(5X,7HKMAX = ,I5,5X,7HLMAX = ,I5)
       454         PRINT 318
       455         PRINT 999
       456   999   FORMAT(5X,15HREGIS4 COMPLETE)
       457         JSTRT=1
       460         JSTOP=20
       461   405   PRINT 403
       462   403   FORMAT(1H1,5X,18HCORRELATION MATRIX)
       463         DO 401 I1=1,NROW1
       464         PRINT 402,(NCOR(I1,J),J=JSTRT,JSTOP)
       471   402   FORMAT(1X,20(I4,1X))
       472   401   CONTINUE
       474         IF(JSTOP.EQ.NCOL1)GO TO 406
       477         JSTRT=JSTOP+1
       500         JSTOP=JSTOP+20
       501         IF(JSTOP.LE.NCOL1)GO TO 405
       504         JSTOP=NCOL1
       505         GO TO 405
       506   406   CONTINUE
       507         STOP
       510         END
```

APPENDIX B. CORRELATION ROUTINE LISTING

```
5245, CAMPBELL     ,140145,00,                    FORTRAN SOURCE LIST
        ISN        SOURCE STATEMENT

          0 $IBFTC COR      DECK
          1       SUBROUTINE COR(NP,NW,NCOR)
            C*****************************************************************************
            C   THIS SUBROUTINE CONTAINS THE STANDARD CORRELATION CALCULATION            *
            C*****************************************************************************
          2       DIMENSION NW(56,56),NP(112,112),NCOR(57,57),FMT(12)
          3       COMMON/NAME1/ NCOL
          4       COMMON/NAME2/ NROW
          5       COMMON/NAME4/ FMT
          6       READ (5,301) (FMT(L),L=1,12)
         13   301 FORMAT(12A6)
         14       DO 10 I1=1,NROW
         15       READ(5,FMT) (NW(I1,K1), K1=1,NCOL)
         22    10 CONTINUE
         24       NROW=NROW+NROW
         25       NCOL=NCOL+NCOL
         26       READ (5,301) (FMT(L),L=1,12)
         33       DO 20 I2=1,NROW
         34       READ(5,FMT) (NP(I2,K2), K2=1,NCOL)
         41    20 CONTINUE
         43       NROW=NROW/2
         44       NCOL=NCOL/2
         45       CALL TIMNOW(IT1)
         46       NROW1=NROW+1
         47       NCOL1=NCOL+1
         50       DO 1 I=1,NROW1
         51       DO 2 J=1,NCOL1
         52       NCOR(I,J)=0
         53       DO 3 K=1,NROW
         54       I1=I+K-1
         55       DO 4 L=1,NCOL
         56       J1=J+L-1
         57       NCOR(I,J)=NCOR(I,J)+NW(K,L)*NP(I1,J1)
         60     4 CONTINUE
         62     3 CONTINUE
         64     2 CONTINUE
         66     1 CONTINUE
         70       CALL TIMNOW(IT2)
         71       JTIME=IT2-IT1
         72       PRINT 318
         73   318 FORMAT(5X,50H**************************************************)
         74       PRINT 317,JTIME
         75   317 FORMAT(/////,5X,15HELAPSED TIME = ,I5,11H/60 SECONDS,///////)
         76       PRINT 318
         77       PRINT 999
        100   999 FORMAT(5X,15HCOR    COMPLETE)
        101       RETURN
        102       END
```

APPENDIX C. FAST FOURIER TRANSFORM CORRELATION LISTING

```
5245, CAMPBELL      ,140145,00,                    FORTRAN SOURCE LIST
      ISN       SOURCE STATEMENT

        0 $IBFTC MAIN
        1       DIMENSION XNWR(64,64),XNWI(64,64),XNPR(64,64),XNPI(54,64)
        2       DIMENSION IA(64,64)
        3       READ 1,NROW,NCOL
        4     1 FORMAT(2I5)
        7       CALL IN(XNPR,XNPI,NROW,NCOL)
       10       CALL IN(XNWR,XNWI,NROW,NCOL)
       11       NTOT=NROW*NCOL
       12       CALL TIMNOW(IT1)
       13       CALL FFT7(XNPR,XNPI,NTOT,NROW,NCOL,1)
       14       CALL FFT7(XNPR,XNPI,NTOT,NROW,NTOT,1)
       15       CALL FFT7(XNWR,XNWI,NTOT,NROW,NCOL,1)
       16       CALL FFT7(XNWR,XNWI,NTOT,NROW,NTOT,1)
       17       XNTOT=NTOT
       20       DO 7 I=1,NROW
       21       DO 8 J=1,NCOL
       22       STORE=XNPR(I,J)*XNWR(I,J)+XNPI(I,J)*XNWI(I,J)
       23       XNPI(I,J)=(XNPI(I,J)*XNWR(I,J)-XNWI(I,J)*XNPR(I,J))/XNTOT
       24       XNPR(I,J)=STORE/XNTOT
       25     8 CONTINUE
       27     7 CONTINUE
       31       CALL FFT7(XNPR,XNPI,NTOT,NROW,NCOL,-1)
       32       CALL FFT7(XNPR,XNPI,NTOT,NROW,NTOT,-1)
       33       CALL TIMNOW(IT2)
       34       ITIME=IT2-IT1
       35       PRINT 14
       36       PRINT 13,ITIME
       37    13 FORMAT(1H1,5X,8HITIME = ,I5,3H/60)
       40       PRINT 14
       41    14 FORMAT(60H************************************************************
             1****)
       42       CALL OUT(XNPR,NROW,NCOL)
       43       STOP
       44       END
```

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